Translating Alloy specifications to the point-free style

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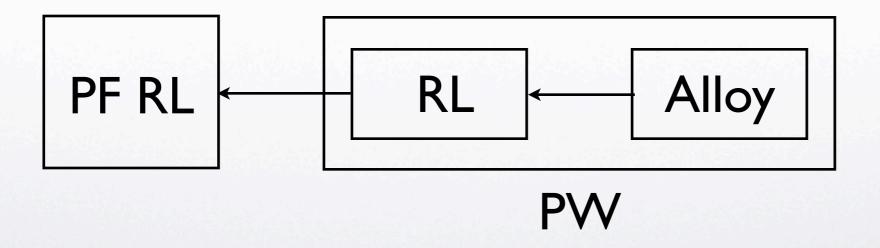
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Motivation

- Alloy provides a tool for automatic bounded verification (the Alloy Analyzer);
- Sometimes however, unbounded verification is necessary;
- Alloy's logic is a kind of relational logic, so relational frameworks are natural choices;
- The point-free (PF) style provides simple enough formulas for manipulation and analysis.

Objectives

 A translation of Alloy models to a PF relational logic (RL) is proposed.



Alloy

- State-based modeling language;
- Simple language, based on simple mathematical notations;
- Characteristics of object modeling;
- Automatic bounded verification.

Calculus of Relations

- Relational Logic (RL):
 - First-order logic enhanced with relational operators (composition, meet, join,...)

$$\langle \forall a, b :: a R b \Rightarrow a S b \rangle$$

- Calculus of Relations (CR):
 - RL without variables

$$R \subseteq S$$

Equivalent to RL with 3 quantified variables.

PF Relational Logic

- Created to overcome the lack of expressiveness of CR;
- Fork Algebras (FA):
 - Introduces a pairs and a new operator fork:

$$c R a \wedge b S a \equiv (c, b) R \nabla S a$$

- Equivalent to RL;
- Categorical CR (CCR):
 - Is able to extend CR and FA with types;
 - Alloy is typed, so it is better suited.

Alloy to FA translation

- Marcelo Frias, Carlos Pombo and Nazareno Aguirre,
 An equational calculus for Alloy;
- Created to translate (only) Alloy formulas to FA;
- Resulting formulas are extremely complex:

all
$$a : A \mid \mathbf{some} \ c : C \mid c \ \mathbf{in} \ r \cdot a$$

$$\overline{\top \cdot \rho(\langle id, \pi_2 \cdot \phi \rangle)} \cdot \langle \pi_1 \cdot \pi_1, \pi_2 \cdot \pi_1, \pi_2, \top \rangle \cdot \langle \pi_1, \pi_2, \top \rangle} \cdot \langle id, \top \rangle \cdot \top = \top$$

$$\phi = id \times \top \cap \overline{\delta(\pi_2 \cdot \pi_1 \cap \pi_2) \cap \overline{\rho(id \times R \cdot \delta(\pi_2 \cdot \pi_1 \cap \pi_2))} \cap id \times \top}$$

Alloy to CCR

- Two-step translation:
 - Alloy to RL: fully expands Alloy formulas to PF
 - RL to CCR: mechanic PW to PF translation, enhanced with heuristics

all a,b : A | (some c : C | c = r·a && c = r·b)
$$\Rightarrow$$
 a = b $\langle \forall a,b \in A :: \langle \exists c \in C :: a \, R \, c \wedge b \, R \, c \rangle \Rightarrow a = b \rangle$

$$\top \subseteq ((\top \cdot (\pi_1 \cdot \pi_2 \cap R \cdot \pi_2) \cap \top \cdot (\pi_2 \cdot \pi_2 \cap R \cdot \pi_2)) \cdot id\nabla \top \cap \overline{\top \cdot (\pi_1 \cap \pi_2)}) \cdot id\nabla \top \cdot \top$$

$$\vdots$$

$$r \cdot r^{\circ} \subseteq id$$

Type System

- Types are represented by the sum of their sub-types;
- Alloy allows the combination of any types with the same arity; all $a : A, b : B \mid a = b$
- Operations are defined to convert two types to their least common super-type.

N-ary relations

 N-ary relations are represented as binary relations with the range as tuples:

$$A \leftarrow B \leftarrow C \rightsquigarrow A \times B \leftarrow C$$

- New operators to manipulate n-ary relations:
 - N-ary composition:

$$(a,b) R \bullet S c \equiv \langle \exists k :: a R k \wedge (k,b) S c \rangle$$

Rotate:

$$(a,b) \overrightarrow{R} c \equiv (c,a) R b$$

Multiplicities

• Signature multiplicities:

Sig A	Property
set	true
some	$\top_{univ \leftarrow univ} \subseteq \top_{univ \leftarrow A} \cdot \top_{A \leftarrow univ}$
lone	$ op _{A\leftarrow A}\subseteq id_{A}$
one	$\mid \top_{univ \leftarrow univ} \subseteq \top_{univ \leftarrow A} \cdot \top_{A \leftarrow univ} \wedge \top_{A \leftarrow A} \subseteq id_A$

Relation multiplicities:

<i>i</i> -th field	Property
set	true
	$\stackrel{n-i}{\rightarrow} \stackrel{n-i}{\rightarrow}$
some	$id\subseteq \overrightarrow{R}\cdot (\overrightarrow{R})^{\circ}$
	$n-i \longrightarrow n-i \longrightarrow$
lone	$(\overrightarrow{R}\)^{\circ}\cdot \overrightarrow{R}\ \subseteq id$
	$\stackrel{n-i}{\rightarrow} \stackrel{n-i}{\rightarrow} \stackrel{n-i}{\rightarrow} \stackrel{n-i}{\rightarrow}$
one	$\mid id \subseteq \overrightarrow{R} \cdot (\overrightarrow{R} \mid)^{\circ} \wedge (\overrightarrow{R} \mid)^{\circ} \cdot \overrightarrow{R} \subseteq id \mid$

Example

Alloy model

```
sig System {
  store: Path \rightarrow lone File,
  open: Handle \rightarrow lone OpenInfo
sig Path {
  dir: one Path
sig Root extends Path { }
sig OpenInfo{
  path: one Path
sig File {}
sig Handle {}
abstract sig Type {}
one sig Reg extends Type {}
one sig Dir extends Type {}
assert ri_ok{
  all s : System | all h: Handle, o: h \cdot (s \cdot open)
               some f: File | f in (o·path)·(s·store)
```

CCR model

```
\begin{aligned} &\textit{Type} \cong \textit{Dir} + \textit{Reg} \\ &\textit{Path} \cong \textit{Path'} + \textit{Root} \\ &\textit{univ} \cong \textit{System} + \textit{Path'} + \textit{Root} + \textit{OpenInfo} + \textit{File} + \textit{Dir} + \textit{Reg} + \textit{Handle} \\ &\textit{Store} :: \textit{System} \times (\textit{Path'} + \textit{Rool}) \leftarrow \textit{File} \\ &\textit{open} :: \textit{System} \times \textit{Handle} \leftarrow \textit{OpenInfo} \\ &\textit{dir} :: (\textit{Path'} + \textit{Root}) \leftarrow (\textit{Path'} + \textit{Root}) \\ &\textit{path} :: \textit{OpenInfo} \leftarrow (\textit{Path'} + \textit{Root}) \end{aligned}
```

Multiplicity

$$store^{\circ} \cdot store \subseteq id$$
 $open^{\circ} \cdot table \subseteq id$ $dir^{\circ} \cdot dir \subseteq id \wedge id \subseteq dir \cdot dir^{\circ}$ $path^{\circ} \cdot path \subseteq id \wedge id \subseteq path \cdot path^{\circ}$

Assert $\langle \forall s, h, o : (s, h) \ table \ o : \langle \exists f, i :: o \ path \ i \land (s, i) \ store \ f \rangle \rangle$

$$\top \cdot (\pi_1 \cap table \cdot \pi_2) \subseteq \top \cdot (\pi_2 \times id \cap path \bullet \overrightarrow{store} \cdot \pi_1 \cdot \pi_1 \cdot \pi_1) \cdot \langle id, \top \rangle$$

Conclusions

- Defined and implemented a highly improved version of the Alloy to PF RL translation;
- Translates not only the formulas but also the signature declarations;
- New operators to deal with n-ary relations, with interesting properties;
- Due to the simplicity, it is suitable for manual or assisted verification.

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